Designing transport-level encryption for datacenter networks

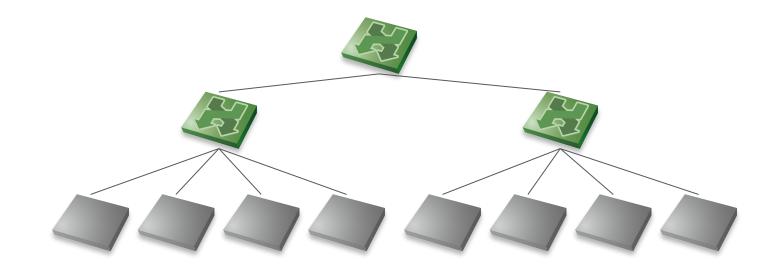
Michio Honda University of Edinburgh



On behalf of my research group members.

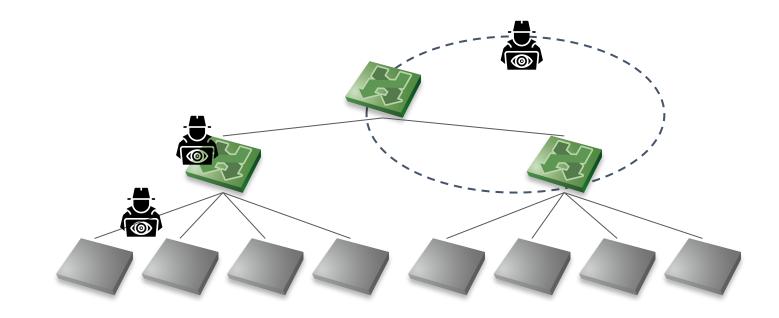
The role of datacenter transport today

- Apps want all of high throughput, low latency & many CPU cycles
- Datacenter transports need good
 - \circ End-to-end congestion control
 - Host stack
 - Switch service



Datacenters need end-to-end encryption

- Multi tenancy
- Multi-vendor hardware/software network components



Datacenter transport requirements

- Modern transport requirements are complex
 - Radically new transport beyond TCP
 - 0-RTT data, receiver-driven congestion control, message boundaries
 - Hardware offload
 - Leaving CPU cycles to the apps
 - In-network compute
 - \circ $\,$ Load balancing, congestion signaling and routing $\,$

Can we design secure datacenter transport without sacrificing those properties?

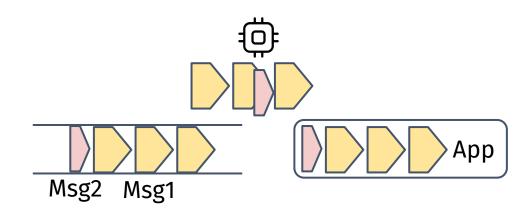
Limitation with bytestream abstractions

• TCP/QUIC

On a packet loss

Head of Line Blocking
Early-arriving small messages should
Msg2 Msg1

On a CPU hotspot



Problems with RDMA abstractions

- Google/Intel Falcon, AWS Scalable Reliable Datagram (SRD)
- Unordered *packet* delivery is supported
 - **NOT** unordered *message delivery*

Design space: Transport-level encryption

	Encrypt.	Abstract.	NIC offload	Wire proto.	Host LB	In-net com.	
TcpCrypt[3]	Inline	Stream	N	TCP	Conn.	N	
QUIC[19]	TLS	Stream	N*	UDP	Conn.	N	*FPGA NIC attempt [52]
TLS/TCP[32]	TLS	Stream	Crypto+TSO	TCP	Conn.	N	
Falcon [8]	PSP	Ordered conns.*	Crypto+TSO	** UDP	Conn.	Ν	*RDMA verbs **Custom NIC or Intel IPU
SDT	TLS	Msg.	Crypto+TSO	New	Msg.	Y*	*With shared key for data muta- tion [46]
Homa[29]/NDF	P[14] -	Msg.	TSO	New	Msg.	Y	
MTP[46]	1.5.	Msg.	TSO	UDP	5	Y	
SRD[43]	-	Dgram.	Full*	Unknown	-	Y	*Custom NIC
KCM[21]/µTCF	·[27] -	Msg.*	TSO	TCP	Conn.	Y*	*high overheads

Table 1: Key properties of encrypted or message-based transports.

Middleground: Unencrypted message-based transport

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Middleground: Unencrypted message-based transport

- Homa*
 - Active development in Linux
 - \circ General to transform to other protocols like NDP
- MTP**
 - Similar to Homa
 - Introduction of in-network compute
 - Load balancing, multipath, congestion signalling, data mutation

Src port	Dst port					
Msg ID						
Msg len						
Msg off						
Payload						

Figure 1: Generalized message-based transport packet format based on MTP [45] and Homa [29]. Shaded parts are identical between the packets that belong to the same message.

> * Ousterhout et al, ATC'21 ** Stephens et al, HotNets'21

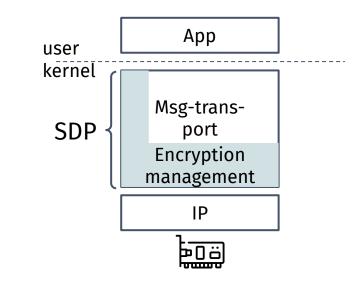
SDP Overview

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C) (2) (2) (3)			140.00		11122-01		

Table 1: Key properties of encrypted or message-based transports.

SDP overview

- Transport-level encryption for datacenter networks
- Message level transport
 - In-network computing support
 - Even data mutation with key sharing
- Opportunistic NIC offload
 - Commodity NVIDIA CX6/7 NICs
- Transport protocol number agnostic
 - \circ Co-existence with existing traffic
- Optional 0-RTT handshake

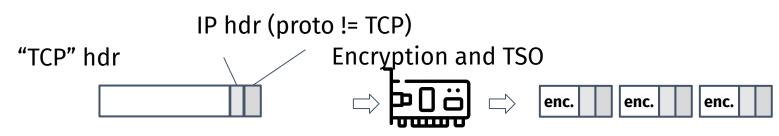


- ~2800 LoC change in Homa/Linux
- ~300 LoC change in the mlx5 driver
- Support Linux 6.2 and 6.6

TLS offload with commodity NICs

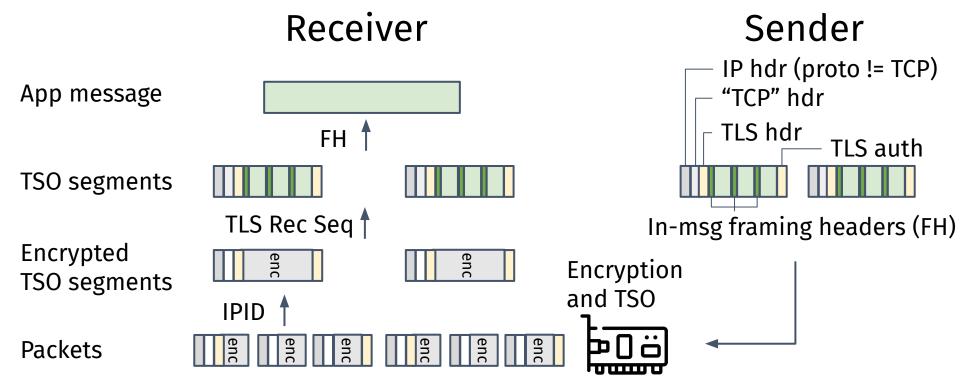
- It is a deal breaker to be able to use existing HW offload
- Full TOE-based approach (Chelsio T6)
 - $\circ~$ Bad even for TCP (e.g., options are gone) and unfavored by operators*
- Autonomous offload* (NVIDIA ConnectX-6/7)
 - Mainstream today
 - Likely similar architecture in Fungible (Microsoft) and Netronome NICs

• It works for non-TCP!



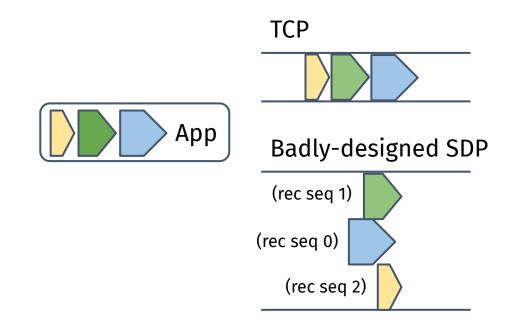
Any-size, unordered authenticated message

- An app message can consist of multiple TSO segments
 - Example below: one app message over two TSO segments
- A TSO segment can consist of multiple packets



Message-level parallelism

- Granularity of parallelism
 - TCP (Connection-level) strict in-order delivery
 - SDP (Message-level) out-order delivery at message level
 - A later message can be received earlier
 - Global record sequence number (over TCP bytestream) no longer works



Result A:

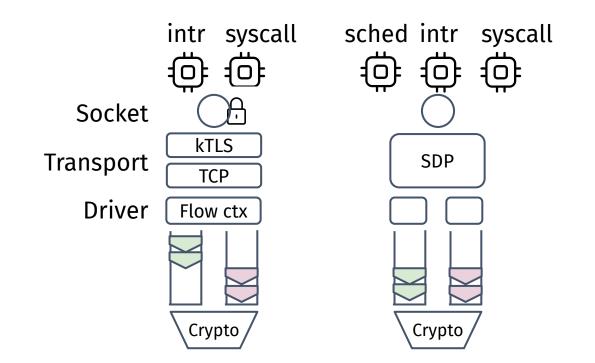
Receiver decrypts record sequence 2 with expecting record sequence 0 -> decrypt failure

Result B:

Receiver waits for record sequence 0 even other records are received -> Head-of-line blocking

NIC offloading

- NICs expect all the data is serialized
 - Under socket lock for TCP
- Message-based transports send multiple messages in parallel in the same flow



Message-level parallelism

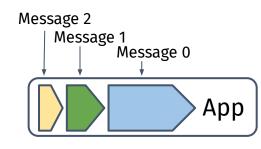
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Solution Assign unique record sequence space to each message

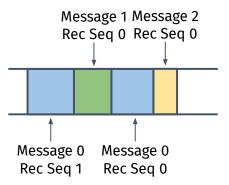


Replay attack protection

- Intra-message: Record sequence numbers increment sequentially like normal TLS
- Inter-Message: Unique message ID used only once in the authenticated session

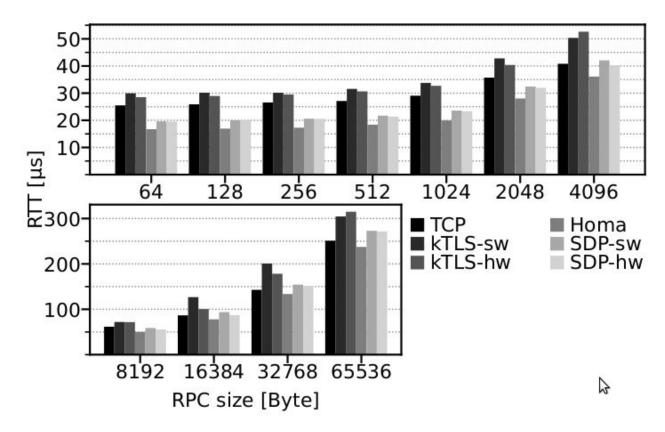


SDP Wire



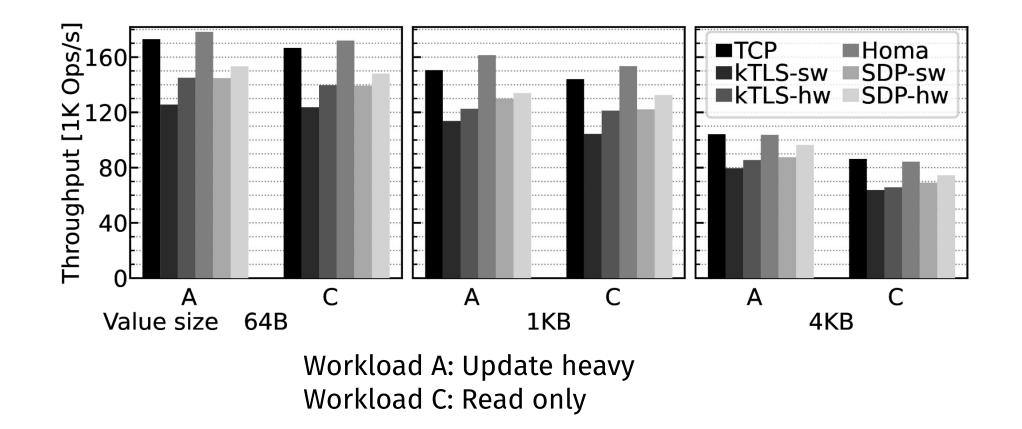
Unloaded latency

- SDP outperforms kTLS by 13–32% with hw offload and 10–35% without it
 - $\circ~$ Homa is faster than TCP by 5–35 %



Redis throughput

• SDP outperforms kTLS by 5–13 % with TLS offload and 8–17 % without it



Summary

- We need security in datacenter networks
- Challenging to preserve important transport properties today:
 - NIC offloading
 - Departure from TCP
 - In-Network Computing support

while preserving the same threat model as TLS/TCP

• SDP solves it

- Existing TLS NIC offload
- Arbitrary-sized, encrypted message
- Same threat model as TLS/TCP
- Protocol number agnostic